## CS61B DISCUSSION 8

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## Revisited: Asymptotics

* Big Oh and Big Omega are useful, but might not give us the best information.
* Example: $\mathrm{n}=\mathrm{O}(\mathrm{n})$. But $\mathrm{n}=\mathrm{O}\left(\mathrm{n}^{\wedge} 2\right)$, and $\mathrm{n}=\mathrm{O}\left(2^{\wedge} \mathrm{n}\right)$ as well! You can give many upper bounds for the same function.
* Better: Use Big Theta. Tighter bound. See Discussion Q1.


## Some summations to know

* $1+2+3+\ldots+\mathrm{N}=\mathrm{N}(\mathrm{N}+1) / 2=\Theta\left(\mathrm{N}^{\wedge} 2\right)$
* $1+2+4+8+\ldots+\mathrm{N}=2 \mathrm{~N}-1 \quad=\Theta(\mathrm{N})$
* You should memorize this, as they will come in handy.
* You don't need to know why. Refer back to your Math 1B notes if you're curious.


## Amortized Analysis

* Way of showing that, on average, what runtime of something is.
* Example: ArrayDeque usually has constant inserts. Every so often, it has to resize, which is in O(n).
* We can show that it Amortized Cost (cost in the long run) is always constant. This is what the table Alan showed in lecture was trying to do.

| Insert \# | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a[i]=$ cost (write cost) | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| resize cost(copy cost) | 0 | 2 | 4 | 0 | 8 | 0 | 0 | 0 | 16 | 0 | 0 | 0 | 0 | 0 |
| total cost | 1 | 3 | 5 | 1 | 9 | 1 | 1 | 1 | 17 | 1 | 1 | 1 | 1 | 1 |
| cumulative cost | 1 | 4 | 9 | 10 | 19 | 20 | 21 | 22 | 39 | 40 | 41 | 42 | 43 | 44 |

## AMORTIZED ANALYSIS: A DIFFERENT APPROACH. (DEMO)

